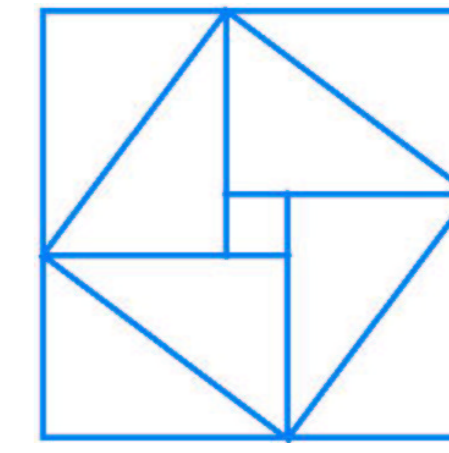


Polynomial general solutions for first order ODEs with constant coefficients

Ruyong Feng and Xiao-Shan Gao
email: {ryfeng, xgao}@mmrc.iss.ac.cn



Institute of Systems Science AMSS
Academia Sinica Beijing 100080, China

1 Basic notations and conceptions

Notation 1.1

\mathbf{K} := differential field of meromorphic functions

$$y_i := \frac{d^i y}{dx^i}$$

$\mathbf{K}\{y\}$:= differential polynomial ring

$F(y)$:= irreducible differential polynomial with constant coefficients

S := separant of $F(y)$

$\Sigma_F := \{A \in \mathbf{K}\{y\} \mid S^k A \equiv 0 \pmod{\{F(y)\}}\}$

$\mathcal{V}_F :=$ differential variety of Σ_F over \mathbf{K}

Definition 1.2 General solutions of $F(y) = 0 \triangleq$ a generic zero of \mathcal{V}_F .

Definition 1.3 Polynomial general solutions of $F(y) = 0 \triangleq$ general solutions of $F(y) = 0$ in the following form

$$\hat{y} = \sum_{i=0}^n a_i x^i, (a_n \neq 0) \quad (1)$$

where a_i are constants.

We will abbreviate Polynomial general solutions to PGSols in the following.

2 Main results

1. A criterion theorem for whether ODEs has PGSols or not.

Theorem 2.1 $F(y) = 0$ has a PGSols iff there is a non-negative integer k such that the pseudo-remainder of y_k wrt $F(y)$ is zero.

Theorem 2.2 Let $F(y)$ be of order o . If $F(y) = 0$ has a PGSols of form (1), then the a_i depend on o independent parameters.

2. An algorithm of complexity $O(n^9)$ to get PGSols of first order ODEs if it exists.

3 PGSols for first order ODEs

In the following, we always assume that $F(y)$ is of order one. From the following theorem, finding PGSols can be reduced to finding a polynomial solution.

Theorem 3.1 Let $\bar{y} = \sum_{i=0}^n \bar{a}_i x^i$ be a solution of $F(y) = 0$, where $\bar{a}_i \in \mathbb{C}$ and $\bar{a}_n \neq 0$. Then

$$\hat{y} = \sum_{i=0}^n \bar{a}_i (x+c)^i \quad (2)$$

is a PGSols for $F(y) = 0$, where c is arbitrary constants.

We will solve the PGSols of $F(y) = 0$ by two steps:

1. Determine the degree bound of PGSols.
2. Find the coefficients of PGSols.

3.1 The degree bound of PGSols

Lemma 3.2 Let $f_1 = \sum_{i=0}^n \bar{a}_i x^i - y$, $f_2 = \sum_{i=1}^n i \bar{a}_i x^{i-1} - y_1$ ($n \geq 1$, $\bar{a}_n \neq 0$, $\bar{a}_i \in \mathbb{C}$). If $n \geq 2$, let R be the Sylvester resultant of f_1 and f_2 wrt x and if $n = 1$, let $R = f_2$. Then R is an irreducible differential polynomial in $\mathbb{C}[y, y_1]$ and has the form

$$R = (-1)^n \bar{a}_n^{n-1} y_1^n + (-1)^{n-1} n^n \bar{a}_n^n y^{n-1} + G(y, y_1) \quad (3)$$

where $\text{tdeg}(G)$ (total degree) $\leq n-1$ and G does not contain the term y^{n-1} .

Theorem 3.3 Use the same notations in Lemma 3.2. If $\bar{y} = \sum_{i=0}^n \bar{a}_i x^i$ is a solution of $F(y) = 0$, then $F(y) = \lambda R$, where $\lambda \in \mathbb{C}$ and $\lambda \neq 0$.

From Lemma 3.2 and Theorem 3.3, if $F(y) = 0$ has PGSols of the form (2), $F(y)$ must be the form

$$F(y) = ay_1^n + by^{n-1} + G(y, y_1) \quad (4)$$

where $a, b \in \mathbb{C}$ and $ab \neq 0$, and G is as in the (3).

Hence the degree of PGSols equals to $\deg(F(y), y_1)$.

3.2 The coefficients of PGSols

Lemma 3.4 Let $F(y)$ be of the form (4) and have a PGSols of the form (2). Then

$$\bar{a}_n = -\frac{b}{n^n a}. \quad (5)$$

Lemma 3.5 Let $F(y) = 0$ have a PGSols of the form (2). Then we may construct a PGSols of the following form for $F(y) = 0$

$$\hat{y} = \bar{a}_n (x+c)^n + \sum_{i=0}^{n-2} \tilde{a}_i (x+c)^i. \quad (6)$$

where \tilde{a}_i are complex numbers.

From above Lemma, we may assume that $\bar{a}_{n-1} = 0$ in PGSols of the form (2).

Lemma 3.6 Let $F(y)$ be of the form (4) and $z = (-b/n^n a)x^n + a_{n-1}x^{n-1} + a_{n-2}x^{n-2} + \dots + a_0$. Substituting y by z in $F(y)$, the coefficients of $x^{(n-1)^2+i-1}$ in $F(z)$ is of the form

$$(-b/n^n a)^{n-2} (n-1-i) b a_i + h_i(a_{n-1}, \dots, a_{i+1}) \quad \text{for } i = n-2, \dots, (7)$$

From Lemma 3.4, 3.5 and 3.6, we may compute the coefficients of PGSols.

4 Algorithm and Example

Algorithm 4.1 The input is $F(y)$. The output is PGSols of $F(y) = 0$ if it exists.

1. If $F(y)$ can be written as the form (4), then goto step 2. Otherwise, $F(y)$ has no PGSols and the algorithm terminates.
2. Let $F(y)$ be of degree n in y_1 . Let $\bar{a}_n = -\frac{b}{n^n a}$, $\bar{a}_{n-1} = 0$, $\bar{a}_i = -\frac{h_i(\bar{a}_{n-1}, \dots, \bar{a}_{i+1})}{((-b/n^n a)^{n-2} (n-1-i) b)}$, $i = n-2, \dots, 0$, where h_i are from Lemma 3.6. We have $\bar{a}_i \in \mathbb{C}$.
3. Let $\bar{y} = \sum_{i=0}^n \bar{a}_i x^i$. If $F(\bar{y}) \equiv 0$ then $\hat{y} = \sum_{i=0}^n \bar{a}_i (x+c)^i$ is PGSols of $F(y) = 0$. Otherwise, $F(y) = 0$ has no PGSols.

Example 4.2 Let

$$F = y_1^4 - 8y_1^3 + (6 + 24y)y_1^2 + 257 + 528y^2 - 256y^3 - 552y$$

1. Write F as

$$F = y_1^4 - 256y^3 + G(y, y_1)$$

where $G = -8y_1^3 + (6 + 24y)y_1^2 + 257 + 528y^2 - 552y$.

2. Let $a_4 = \frac{256}{4^4} = 1$ and $\bar{y} = x^4 + a_2 x^2 + a_1 x + a_0$. Replacing y by \bar{y} in F , the coefficients of x^{10}, x^9, x^8 are

$$\begin{aligned} &384 - 256a_2 \\ &-512 - 512a_1 \\ &768a_2 + 528 - 384a_2^2 - 768a_0. \end{aligned}$$

Then we have $a_2 = \frac{3}{2}$, $a_1 = -1$, $a_0 = \frac{17}{16}$

3. Let $\bar{y} = x^4 + \frac{3}{2}x^2 - x + \frac{17}{16}$. Substituting y by \bar{y} into F , F becomes zero. Hence PGSols of $F = 0$ is

$$\hat{y} = (x+c)^4 + \frac{3}{2}(x+c)^2 - (x+c) + \frac{17}{16}$$